

#### Campus Networking Workshop

Layer-2 Network Design







# Layer 2 Concepts

- Layer 2 protocols basically control access to a shared medium (copper, fiber, electromagnetic waves)
- Ethernet is the *de-facto* standard today
  - -Reasons:
    - Simple
    - Cheap
    - Manufacturers keep making it faster





#### Ethernet Functions

- Source and Destination identification
  - -MAC addresses
- Detect and avoid frame collisions
  - -Listen and wait for channel to be available
  - If collision occurs, wait a random period before retrying
    - This is called CASMA-CD: Carrier Sense Multiple Access with Collision Detection
  - -1Gbps links and above are always full duplex





# Evolution of Ethernet Topologies

• Bus

-Everybody on the same coaxial cable

- Star
  - -One central device connects every other node
    - First with hubs (repeated traffic)
    - Later with switches (bridged traffic)
  - -Structured cabling for star topologies standardized





# •Switched Star Topology Benefits

- It's modular:
  - -Independent wires for each end node
  - -Independent traffic in each wire
  - A second layer of switches can be added to build a hierarchical network that extends the same two benefits above
  - -ALWAYS DESIGN WITH MODULARITY IN MIND





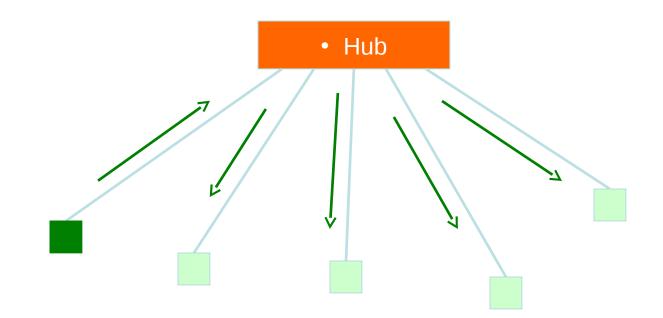
#### •Hub

- Receives a frame on one port and sends it out <u>every other port, always</u>.
- Collision domain spans the whole hub or chain of hubs
- Traffic ends up in places where it's not needed









- A frame sent by one node is always sent to every other node.
- Hubs are also called "repeaters" because they just "repeat"
- what they hear.





#### Switch

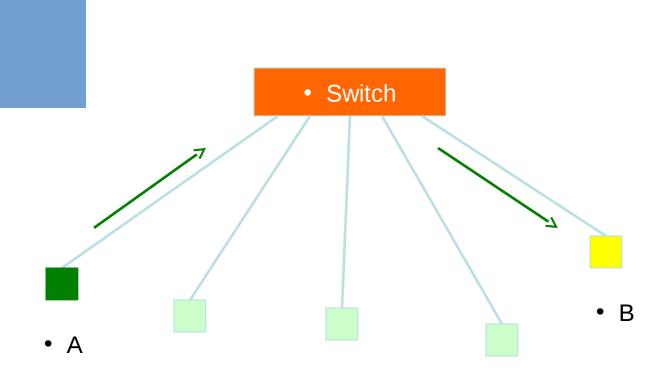
- *Learns* the location of each node by looking at the source address of each incoming frame, and builds a *forwarding table*
- **Forwards** each incoming frame only to the port where the destination node is
  - Reduces the collision domain
  - Makes more efficient use of the wire
  - Nodes don't waste time checking frames not destined to them





#### •Switch

- Forwarding Table
- Address
   Port
- AAAAAAAAA 1 AA
- BBBBBBBBBB 5 BB







#### Switches and Broadcast

- A switch broadcasts some frames:
  - –When the destination address is not found in the table
  - -When the frame is destined to the broadcast address (FF:FF:FF:FF:FF:FF)
  - –When the frame is destined to a multicast ethernet address
- So, switches do not reduce the broadcast domain!





#### Switch vs. Router

- Routers more or less do with IP packets what switches do with Ethernet frames
  - A router looks at the IP packet destination and checks its *forwarding table* to decide where to forward the packet
- Some differences:
  - -IP packets travel inside ethernet frames
  - -IP networks can be logically segmented into *subnets*
  - Switches do not usually know about IP, they only deal with Ethernet frames





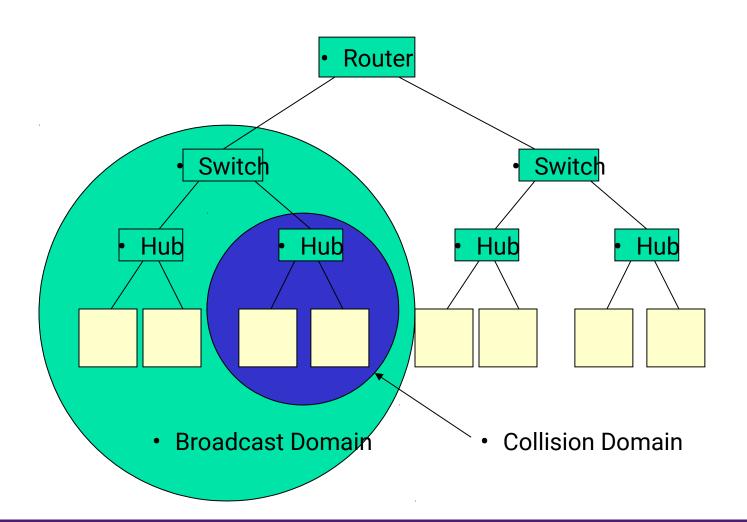
#### Switch vs. Router

- Routers do not forward Ethernet broadcasts. So:
  - -Switches reduce the <u>collision domain</u> -Routers reduce the <u>broadcast domain</u>
- This becomes *really* important when trying to design hierarchical, scalable networks that can grow sustainably





#### Traffic Domains





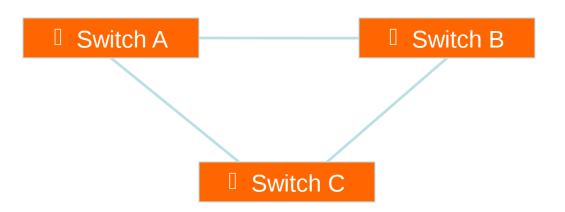


### •Traffic Domains

- Try to eliminate collision domains –Get rid of hubs!
- Try to keep your broadcast domain limited to no more than 250 simultaneously connected hosts
  - -Segment your network using routers







- When there is more than one path between two switches
- What are the potential problems?

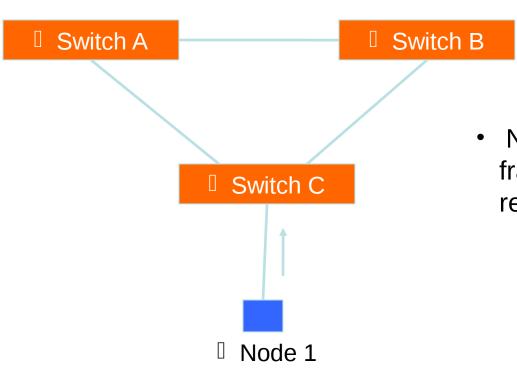




- If there is more than one path between two switches:
  - -Forwarding tables become unstable
    - Source MAC addresses are repeatedly seen coming from different ports
  - -Switches will broadcast each other's broadcasts
    - All available bandwidth is utilized
    - Switch processors cannot handle the load



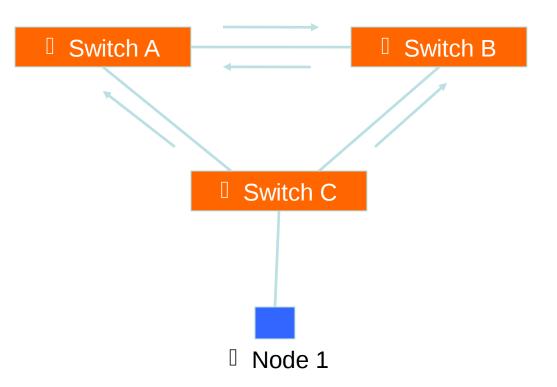




 Node1 sends a broadcast frame (e.g. an ARP request)



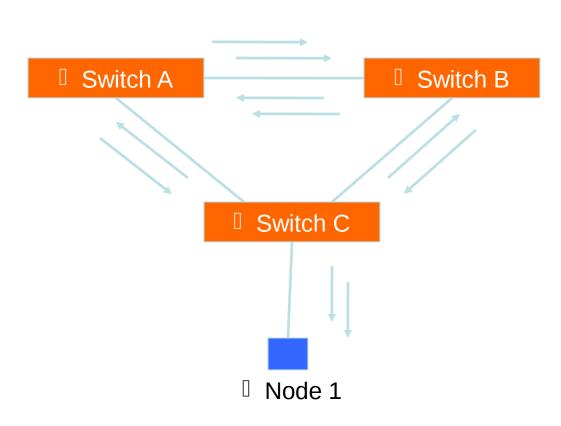




• Switches A, B and C broadcast node 1's frame out every port







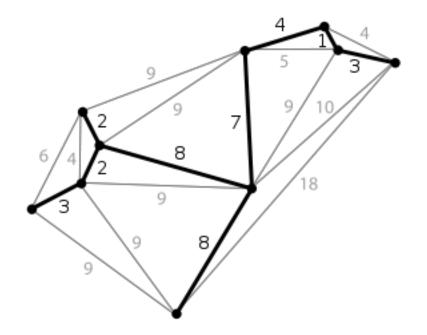
- But they receive each other's broadcasts, which they need to forward again out every port!
- The broadcasts are amplified, creating a broadcast storm



- But you can take advantage of loops! –Redundant paths improve resilience when:
  - A switch fails
  - Wiring breaks
- How to achieve redundancy without creating dangerous traffic loops?

# What is a Spanning Tree

- "Given a connected, undirected graph, a *spanning tree* of that graph is a subgraph which is a tree and connects all the vertices together".
- A single graph can have many different spanning trees.





• The purpose of the protocol is to have bridges dynamically discover a subset of the topology that is loop-free (a tree) and yet has just enough connectivity so that where physically possible, there is a path between every switch



- Several flavors:
  - -Traditional Spanning Tree (802.1d)
  - -Rapid Spanning Tree or RSTP (802.1w)
  - -Multiple Spanning Tree or MSTP (802.1s)



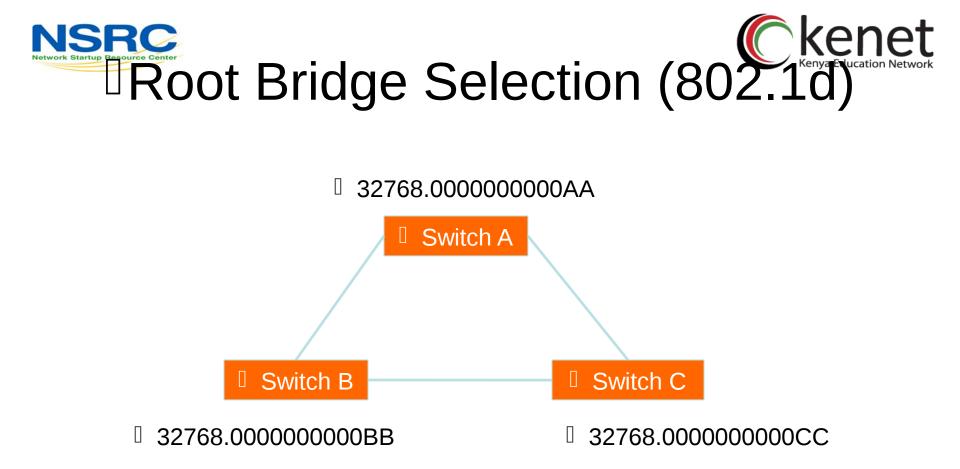
- Switches exchange messages that allow them to compute the Spanning Tree
  - -These messages are called BPDUs (Bridge Protocol Data Units)
  - -Two types of BPDUs:
    - Configuration
    - Topology Change Notification (TCN)



- First Step:
  - -Decide on a point of reference: the *Root Bridge*
  - –The election process is based on the Bridge ID, which is composed of:
    - <u>The Bridge Priority</u>: A two-byte value that is configurable
    - <u>The MAC address</u>: A unique, hardcoded address that cannot be changed.



- Each switch starts by sending out BPDUs with a Root Bridge ID equal to its own Bridge ID –*I am the root!*
- Received BPDUs are analyzed to see if a <u>lower</u> Root Bridge ID is being announced
  - If so, each switch replaces the value of the advertised Root Bridge ID with this new lower ID
- Eventually, they all agree on who the Root Bridge is



- All switches have the same priority.
- Who is the elected root bridge?



- Now each switch needs to figure out where it is in relation to the Root Bridge
  - -Each switch needs to determine its *Root Port*
  - -The key is to find the port with the lowest **Root Path Cost** 
    - The cumulative cost of all the links leading to the Root Bridge



 Root Path Cost is the accumulation of a link's Path Cost and the Path Costs learned from neighboring Switches.
 It answers the question: How much does it cost to reach the Root Bridge through this port?



1.Root Bridge sends out BPDUs with a Root Path Cost value of 0

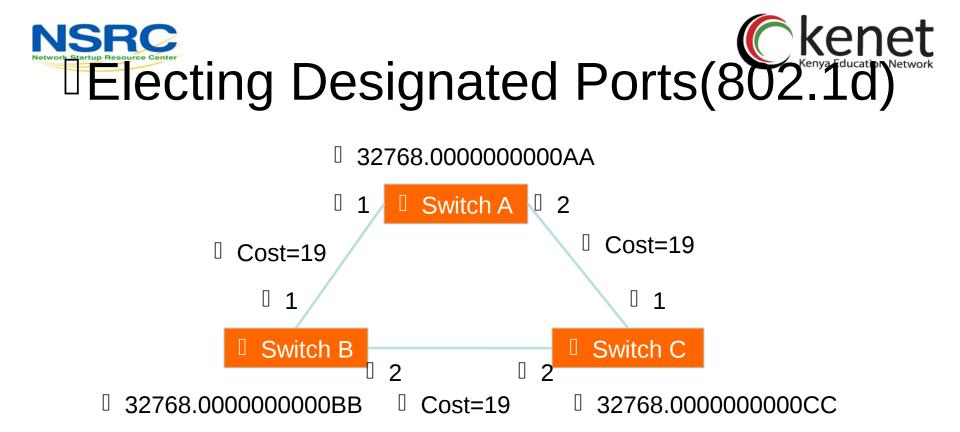
- 2.Neighbor receives BPDU and adds port's Path Cost to Root Path Cost received
- 3.Neighbor sends out BPDUs with new cumulative value as Root Path Cost4.Other neighbors down the line keep adding in the same fashion



- On each switch, the port where the lowest Root Path Cost was received becomes the *Root Port*
  - This is the port with the best path to the Root Bridge

# Electing Designated Potskenet (802.1d)

- OK, we now have selected root ports but we haven't solved the loop problem yet, have we – The links are still active!
- Each network segment needs to have only one switch forwarding traffic to and from that segment
- Switches then need to identify one *Designated Port* per link
  - The one with the lowest cumulative Root Path Cost to the Root Bridge



• Which port should be the Designated Port on each segment?

# Electing Designated Potskenet (802.1d)

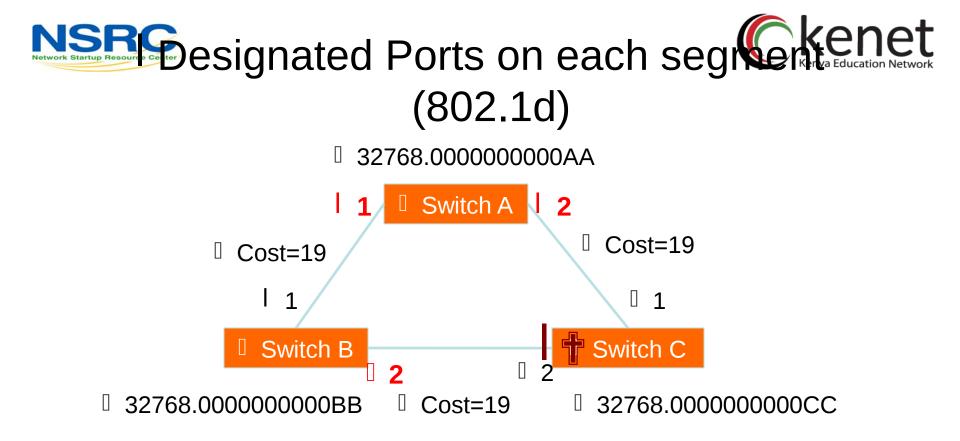
- Two or more ports in a segment having identical Root Path Costs is possible, which results in a tie condition
- All STP decisions are based on the following sequence of conditions:
  - -Lowest Root Bridge ID
  - -Lowest Root Path Cost to Root Bridge
  - -Lowest Sender Bridge ID
  - -Lowest Sender Port ID





# Blocking a port

- Any port that is not elected as either a Root Port, nor a Designated Port is put into the Blocking State.
- This step effectively breaks the loop and completes the Spanning Tree.



 Port 2 in Switch C is then put into the *Blocking State* because it is neither a Root Port nor a Designated Port



- Disabled
  - -Port is shut down
- Blocking
  - -Not forwarding frames
  - -Receiving BPDUs
- Listening
  - -Not forwarding frames
  - -Sending and receiving BPDUs



- Learning
  - -Not forwarding frames
  - -Sending and receiving BPDUs
  - -Learning new MAC addresses

#### • Forwarding

- -Forwarding frames
- -Sending and receiving BPDUs
- -Learning new MAC addresses



- Switches will recalculate if:
  - -A new switch is introduced
    - It could be the new Root Bridge!
  - –A switch fails
  - –A link fails



- Using default STP parameters might result in an undesired situation
  - -Traffic will flow in non-optimal ways
  - An unstable or slow switch might become the root
- You need to plan your assignment of bridge priorities carefully



- Some vendors have included features that protect the STP topology:
  - -Root Guard
  - -BPDU Guard
  - -Loop Guard
  - -UDLD
  - -Etc.



- Enable spanning tree even if you don't have redundant paths
- Always plan and set bridge priorities —Make the root choice deterministic —Include an alternative root bridge
- If possible, do not accept BPDUs on end user ports
  - -Apply BPDU Guard or similar where available



- Moving from the Blocking state to the Forwarding State takes at least 2 x *Forward Delay* time units (~ 30 secs.)
   This can be annoying when connecting end user stations
- Some vendors have added enhancements such as *PortFast*, which will reduce this time to a minimum for edge ports
  - Never use PortFast or similar in switch-to-switch links
- Topology changes typically take 30 seconds too — This can be unacceptable in a production network



- Convergence is much faster

   Communication between switches is more
   interactive
- Edge ports don't participate
  - Edge ports transition to forwarding state immediately
  - –If BPDUs are received on an edge port, it becomes a non-edge port to prevent loops

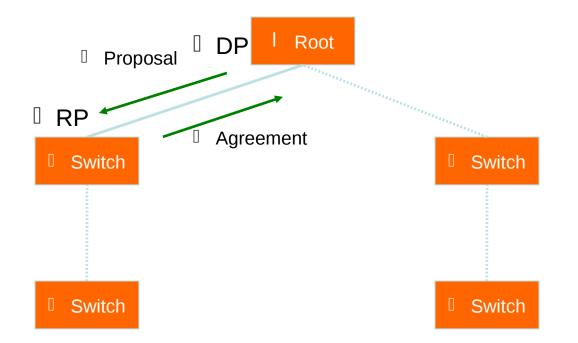
# Rapid Spanning Tree (802.1W)

- Defines these port roles:
  - -Root Port (same as with 802.1d)
  - -Alternate Port
    - A port with an alternate path to the root
  - –Designated Port (same as with 802.1d)
  - -Backup Port
    - A backup/redundant path to a segment where another bridge port already connects.

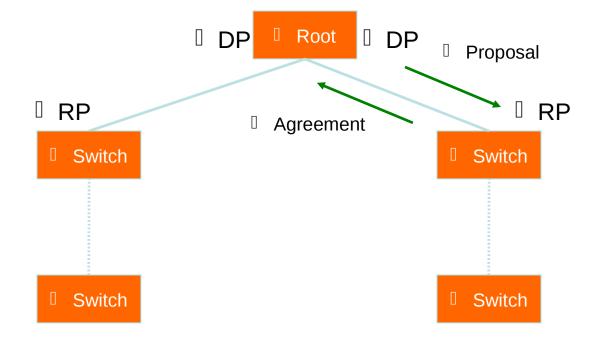


- Synchronization process uses a handshake method
  - After a root is elected, the topology is built in cascade, where each switch proposes to be the designated bridge for each point-to-point link
    While this happens, all the downstream switch
    - links are blocking

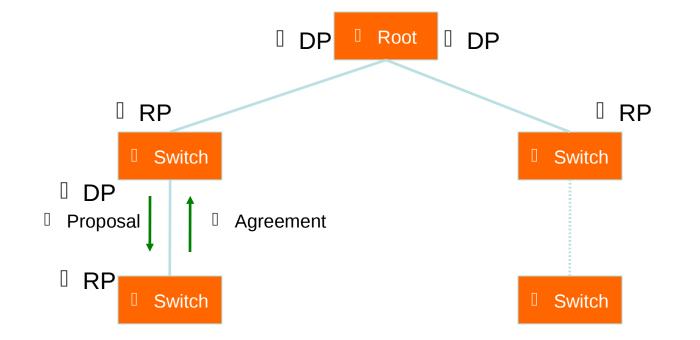
# Rapid Spanning Tree (802.1W)

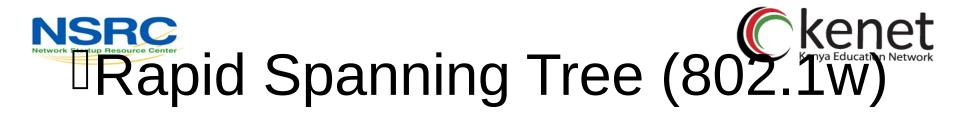


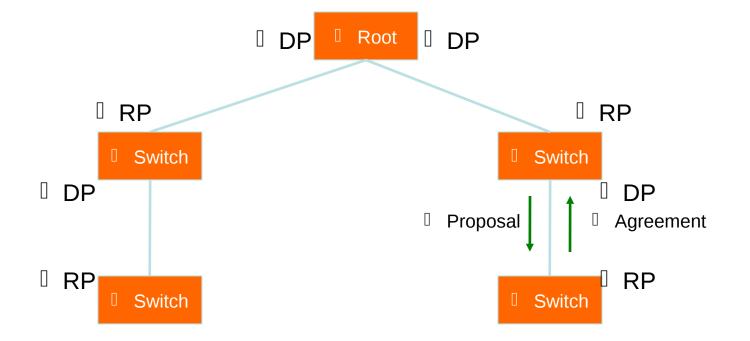














- Prefer RSTP over STP if you want faster convergence
- Always define which ports are edge ports



#### Questions?





- Allow us to split switches into separate (virtual) switches
- Only members of a VLAN can see that VLAN's traffic
  - -Inter-vlan traffic must go through a router
- Allow us to reuse router interfaces to carry traffic for separate subnets

   –E.g. sub-interfaces in Cisco routers





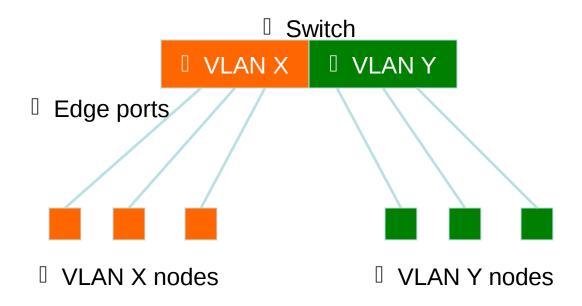


- 2 VLANs or more within a single switch
- *Edge ports*, where end nodes are connected, are configured as members of a VLAN
- The switch behaves as several virtual switches, sending traffic only within VLAN members



#### **Local VLANs**







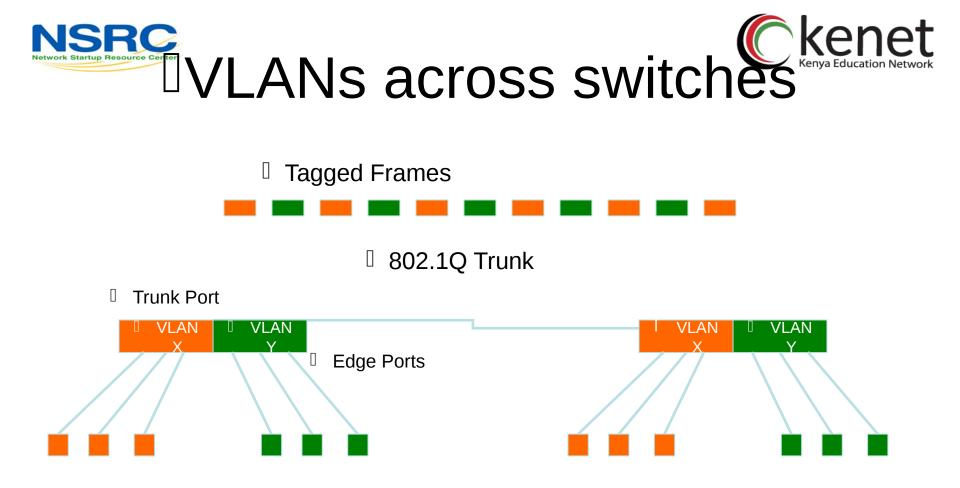
- Two switches can exchange traffic from one or more VLANs
- Inter-switch links are configured as *trunks*, carrying frames from all or a subset of a switch's VLANs
- Each frame carries a *tag* that identifies which VLAN it belongs to







- The IEEE standard that defines how ethernet frames should be *tagged* when moving across switch trunks
- This means that switches from *different vendors* are able to exchange VLAN traffic.



#### <sup>[]</sup> This is called "VLAN Trunking"



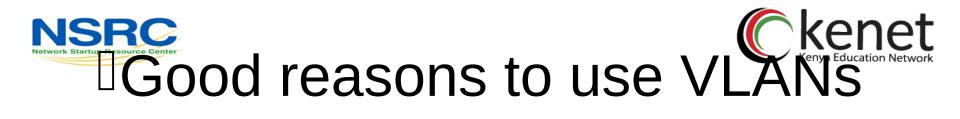


# Tagged vs. Untagged

- Edge ports are not tagged, they are just "members" of a VLAN
- You only need to tag frames in switch-toswitch links (trunks), when transporting multiple VLANs
- A trunk can transport both tagged and untagged VLANs
  - As long as the two switches agree on how to handle those



- You can no longer "just replace" a switch –Now you have VLAN configuration to maintain and backup Eiglel to charie in a constant of the second seco
  - -Field technicians need more skills
- You have to make sure that all the switchto-switch trunks are carrying all the necessary VLANs
  - Need to keep in mind when adding/removing
     VLANs



- You want to segment your network into multiple subnets, but can't buy enough switches
  - -Hide sensitive infrastructure like IP phones, building controls, etc.
- Separate control traffic from user traffic –Restrict who can access your switch management address



- Because you can, and you feel cool  $\ensuremath{\mathbb{Z}}$
- Because they will completely secure your hosts (or so you think)
- Because they allow you to extend the same IP network over multiple separate buildings
  - -This is actually very common, but a bad idea



- Extending a VLAN to multiple buildings across trunk ports
- Bad idea because:
  - -Broadcast traffic is carried across all trunks from one end of the network to another
  - -Broadcast storm can spread across the extent of the VLAN, and affect all VLANS!
  - -<u>Maintenance and troubleshooting nightmare</u>





# **Link Aggregation**

- Also known as port bundling, link bundling
- You can use multiple links in parallel as a single, logical link
  - -For increased capacity
  - -For redundancy (fault tolerance)
- LACP (Link Aggregation Control Protocol) is a standardized method of negotiating these bundled links between switches



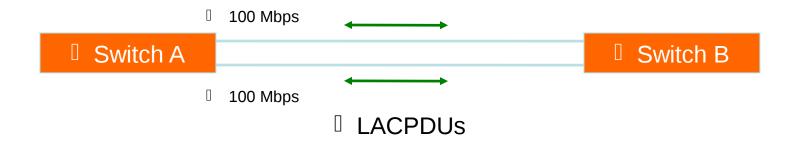


# **LACP** Operation





#### **LACP** Operation

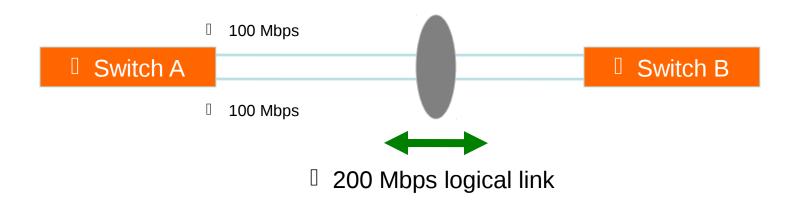


- Switches A and B are connected to each other using two sets of Fast Ethernet ports
- LACP is enabled and the ports are turned on
- Switches start sending LACPDUs, then negotiate how to set up the aggregation





#### **LACP** Operation



- The result is an aggregated 200 Mbps logical link
- The link is also fault tolerant: If one of the member links fail, LACP will automatically take that link off the bundle, and keep sending traffic over the remaining link

### **NDIST**ibuting Traffic In Bundled Links



- Bundled links distribute frames using a hashing algorithm, based on:
  - Source and/or Destination MAC address
  - Source and/or Destination IP address
  - Source and/or Destination Port numbers
- This can lead to unbalanced use of the links, depending on the nature of the traffic
- Always choose the load-balancing method that provides the most distribution



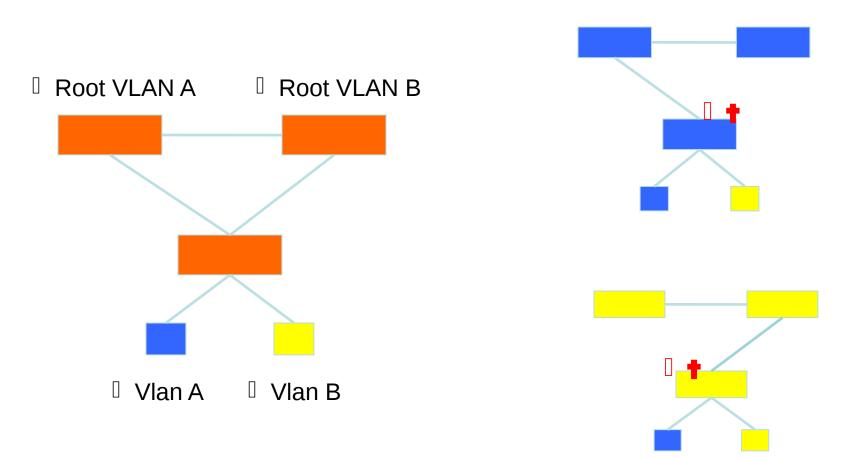
#### Questions?



# Multiple Spanning Treekenet (802.1s)

- Allows separate spanning trees per VLAN group
  - Different topologies allow for load balancing between links
  - -Each group of VLANs are assigned to an "instance" of MST
- Compatible with STP and RSTP

### Multiple Spanning Treekenet (802.1s)



# Multiple Spanning Treekenet (802.1s)

- MST Region
  - -Switches are members of a region if they have the same set of attributes:
    - MST configuration name
    - MST configuration revision
    - Instance-to-VLAN mapping
  - -A digest of these attributes is sent inside the BPDUs for fast comparison by the switches
  - One region is usually sufficient

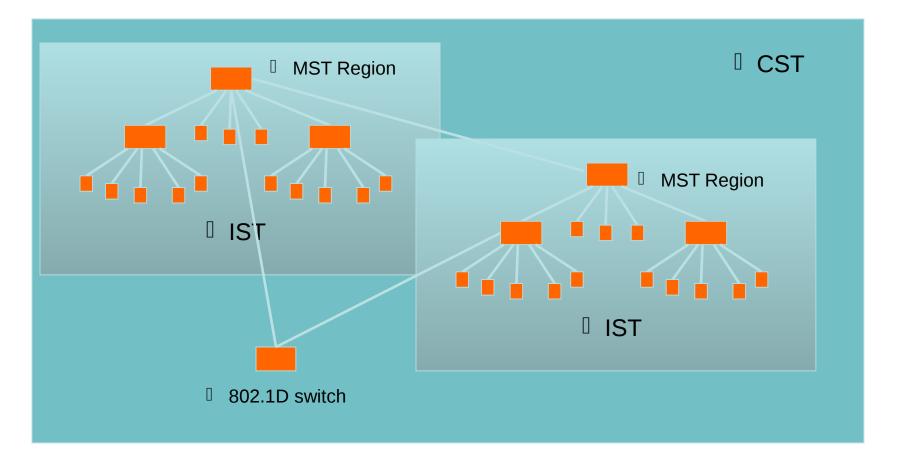
 CST = Common Spanning Tree

 In order to interoperate with other versions of Spanning Tree, MST needs a common tree that contains all the other islands, including other

MST regions

- IST = Internal Spanning Tree
  - -Internal to the Region, that is
  - -Presents the entire region as a single virtual bridge to the CST outside

- MST Instances
  - -Groups of VLANs are mapped to particular Spanning Tree instances
  - -These instances will represent the alternative topologies, or forwarding paths
  - -You specify a root and alternate root for each instance



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- Design Guidelines
  - Determine relevant forwarding paths, and distribute your VLANs equally into instances matching these topologies
  - Assign different root and alternate root switches to each instance
  - -Make sure all switches match region attributes
  - Do not assign VLANs to instance 0, as this is used by the IST





- Minimum features:
  - -Standards compliance
  - -Encrypted management (SSH/HTTPS)
  - -VLAN trunking
  - -Spanning Tree (RSTP at least)
  - -SNMP
    - At least v2 (v3 has better security)
    - Traps





- Other recommended features: –DHCP Snooping
  - Prevent end-users from running a rogue DHCP server
    - Happens a lot with little wireless routers (Netgear, Linksys, etc) plugged in backwards
  - Uplink ports towards the legitimate DHCP server are defined as "trusted". If DHCPOFFERs are seen coming from any untrusted port, they are dropped.





- Other recommended features:
  - -Dynamic ARP inspection
    - A malicious host can perform a man-in-the-middle attack by sending gratuitous ARP responses, or responding to requests with bogus information
    - Switches can look inside ARP packets and discard gratuitous and invalid ARP packets.





- Other recommended features:
  - -IGMP Snooping:
    - Switches normally flood multicast frames out every port
    - Snooping on IGMP traffic, the switch can learn which stations are members of a multicast group, thus forwarding multicast frames only out necessary ports
    - Very important when users run Norton Ghost, for example.



- Enable SNMP traps and/or syslog
  - -Collect and process in centralized log server
    - Spanning Tree Changes
    - Duplex mismatches
    - Wiring problems
- Monitor configurations
  - -Use RANCID to report any changes in the switch configuration



- Collect forwarding tables with SNMP
  - Allows you to find a MAC address in your network quickly
  - You can use simple text files + grep, or a web tool with DB backend
- Enable LLDP (or CDP or similar)
  - -Shows how switches are connected to each other and to other network devices





#### Documentation

- Document where your switches are located
  - -Name switch after building name
    - E.g. building1-sw1
  - -Keep files with physical location
    - Floor, closet number, etc.
- Document your edge port connections
  - –Room number, jack number, server name





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# Thank You

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